

/ BYZANTINE AGREEMENT PROTOCOL */*
/ this files contains the proof of agreement */*

/-----ASSUMPTIONS-----*/*

/ 1-4: there cannot be M votes for v and M votes for v' */*
/ these follows from the fact if this was not true there would be */*
/ $M + M = (N - 2T) + (N - 2T) = N - T + (N - 3T) > N - T$ honest parties */*
/ which is a contradiction since there are only N-T honest parties */*

```
forall (r in ROUNDS) {
  assumption1[r] : assert G  $\neg$ ( main_votes[r][2]=M  $\wedge$  main_votes[r][0]=M);
  assumption2[r] : assert G  $\neg$ ( main_votes[r][2]=M  $\wedge$  main_votes[r][1]=M);
  assumption3[r] : assert G  $\neg$ ( main_votes[r][0]=M  $\wedge$  main_votes[r][1]=M);
  assume assumption1[r], assumption2[r], assumption3[r];
}
forall (r in ROUNDS) for(c = 0; c  $\leq$  1; c = c + 1) {
  assumption4[r][c] : assert G  $\neg$ ( pre_votes[r][c][0]=M  $\wedge$  pre_votes[r][c][1]=M);
  assume assumption4[r][c];
}
```

/-----ADDITIONAL INVARIANTS-----*/*

/ these are proved in a separate file (lemmas.smv) */*

```
forall (r in ROUNDS) forall (n in VOTES) {
  lemma5[r][n] : assert G ( pre[r]=n  $\Rightarrow$  G (pre[r] $\geq$ n) );
  assume lemma5[r][n];
}
forall (r in ROUNDS) for(c = 0; c  $\leq$  1; c = c + 1) for(v = 0; v  $\leq$  1; v = v + 1) forall (n in VOTES) {
  lemma9[r][c][v][n] : assert G ( pre_votes[r][c][v] $\geq$ n  $\Rightarrow$  G (pre_votes[r][c][v] $\geq$ n) );
  assume lemma9[r][c][v][n];
}
forall (r in ROUNDS) for(c = 0; c  $\leq$  1; c = c + 1) for(v = 0; v  $\leq$  1; v = v + 1) {
  lemma11[r][c][v] : assert G (pre_votes[r][c][v]>0  $\Rightarrow$  G (pre_votes[r][c][v]>0) );
  assume lemma11[r][c][v];
}
forall (r in ROUNDS) for(c = 0; c  $\leq$  1; c = c + 1) {
  lemma16[r][c] : assert G ( pre_votes[r][c][1]=0  $\Rightarrow$  pre_votes[r][c][0]=pre[r] );
  assume lemma16[r][c];
}
forall (r in ROUNDS) for(c = 0; c  $\leq$  1; c = c + 1) {
  lemma17[r][c] : assert G ( pre_votes[r][c][0]=0  $\Rightarrow$  pre_votes[r][c][1]=pre[r] );
  assume lemma17[r][c];
}
forall (r in ROUNDS) forall (n in VOTES) {
  lemma19[r][n] : assert G ( main[r] $\geq$ n  $\Rightarrow$  G (main[r] $\geq$ n) );
  assume lemma19[r][n];
}
forall (r in ROUNDS) for(v = 0; v  $\leq$  2; v = v + 1) forall (n in VOTES) {
  lemma23[r][v][n] : assert G ( main_votes[r][v] $\geq$ n  $\Rightarrow$  G (main_votes[r][v] $\geq$ n) );
  assume lemma23[r][v][n];
}
forall (r in ROUNDS) for(v = 0; v  $\leq$  2; v = v + 1) {
  lemma25[r][v] : assert G ( main_votes[r][v]>0  $\Rightarrow$  G (main_votes[r][v]>0) );
  assume lemma25[r][v];
}
forall (r in ROUNDS) {
  lemma31[r] : assert G ( (main_votes[r][1]=0  $\wedge$  main_votes[r][2]=0)  $\Rightarrow$  main_votes[r][0]=main[r] );
}
```

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    assume lemma31[r];
  }
forall (r in ROUNDS) {
  lemma32[r] : assert G ( (main_votes[r][0]=0 ∧ main_votes[r][2]=0) ⇒ main_votes[r][1]=main[r] );
  assume lemma32[r];
}
forall (r in ROUNDS) {
  lemma33[r] : assert G ( (main_votes[r][0]=0 ∧ main_votes[r][1]=0) ⇒ main_votes[r][2]=main[r] );
  assume lemma33[r];
}
forall (r in ROUNDS) forall (i in PROC) {
  lemma37[r][i] : assert G ( decide[r][i] ⇒ G (decide[r][i]) );
  assume lemma37[r][i];
}
forall (r in ROUNDS) for (v = 0; v ≤ 2; v = v + 1) {
  corrupted1[r][v] : assert G ( corrupted_main[r][v] ⇒ G (corrupted_main[r][v]) );
  assume corrupted1[r][v];
}
forall (r in ROUNDS) for (c = 0; c ≤ 1; c = c + 1) for (v = 0; v ≤ 1; v = v + 1) {
  corrupted2[r][c][v] : assert G ( corrupted_pre[r][c][v] ⇒ G (corrupted_pre[r][c][v]) );
  assume corrupted2[r][c][v];
}
forall (r in ROUNDS) for (v = 0; v ≤ 2; v = v + 1) {
  corrupted5[r][v] : assert G ( main_votes[r][v]>0 ⇒ corrupted_main[r][v] );
  assume corrupted5[r][v];
}
forall (r in ROUNDS) for (c = 0; c ≤ 1; c = c + 1) for (v = 0; v ≤ 1; v = v + 1) {
  corrupted6[r][c][v] : assert G ( pre_votes[r][c][v]>0 ⇒ corrupted_pre[r][c][v] );
  assume corrupted6[r][c][v];
}
forall (r in ROUNDS) for(c = 0; c ≤ 1; c = c + 1) {
  coin2[r][c] : assert G ( coin[r]=c ⇒ G ( coin[r]=c ) );
  assume coin2[r][c];
}
forall (r in ROUNDS) {
  lemma40[r] : assert G ( (main_votes[r][0]>0 ∨ corrupted_main[r][0]) ⇒ ( main_votes[r][1]=0 ∧ ¬corrupted_main[r][1] ) );
  assume lemma40[r];
}
forall (r in ROUNDS) {
  lemma41[r] : assert G ( (main_votes[r][1]>0 ∨ corrupted_main[r][1]) ⇒ ( main_votes[r][0]=0 ∧ ¬corrupted_main[r][0] ) );
  assume lemma41[r];
}
}

```

/-----SUBLEMMAS FOR AGREEMENT-----*/*

/ if a party decides on v then there are no main votes for !v and less than M for abstain */*

```

forall (i in PROC) forall (r in ROUNDS) for(v = 0; v ≤ 1; v = v + 1) {
  inv1[i][r][v] : assert G ( (decide[r][i] ∧ decide_value[r][i]=v) ⇒
    (main_votes[r][v]=M ∧ main_votes[r][v+1 mod 2]=0 ∧ ¬corrupted_main[r][v+1 mod 2] ∧ main_votes[r][2]<M) );
  using lemma40[r],
    lemma41[r],
    lemma23[r][v][M],
    corrupted1[r],
    /* assumptions */
    assumption1[r],
    assumption2[r],
    /* abstractions */
    VOTES⇒{0,M},
    /* free variables */
    corrupted_main//free,

```

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    corrupted_main[r][v+1 mod 2],
    corrupted_pre//free,
    main//free,
    main_votes//free,
    main_votes[r][v],
    pre//free,
    pre_votes//free,
    pre_proc//free,
    pre_proc_votes//free
    prove inv1[i][r][v];
}
/* if a party decides on v then no party has a pre_vote for !v in the next round */
forall (i in PROC) forall (r in ROUNDS) for(v = 0; v ≤ 1; v = v + 1) for (c = 0; c ≤ 1; c = c + 1){
  inv2[i][r][v][c] : assert G ( (decide[r][i] ∧ decide_value[r][i]=v) ⇒
    (pre_votes[r+1][c][v+1 mod 2]=0 ∧ ¬corrupted_pre[r+1][c][v+1 mod 2] ) );
  forall (j in PROC) {
    subcase inv2[i][r][v][c][j] of inv2[i][r][v][c] for j=history_pre_votes[r+1][c][v+1 mod 2];
    using inv1[i][r][v],
      lemma19[r][M],
      lemma23[r][0][M],
      lemma23[r][1][M],
      lemma23[r][2][M],
      lemma25[r][v+1 mod 2],
      lemma25[r][2],
      lemma31[r],
      lemma32[r],
      lemma33[r],
      corrupted1[r],
      /* abstractions */
      VOTES⇒{0,M},
      ROUNDS⇒{0,r,r+1},
      /* free variables */
      coin//free,
      corrupted_main//free,
      corrupted_pre//free,
      corrupted_pre[r+1][c][v+1 mod 2],
      decide//free,
      decide[r][i],
      main//free,
      main_votes//free,
      pc//free,
      pc[j],
      pre//free,
      pre_proc//free,
      pre_proc_votes//free,
      pre_votes//free,
      pre_votes[r+1][c][v+1 mod 2],
      round//free,
      round[j]
      prove inv2[i][r][v][c][j];
  }
}
/* if a party decides on 0, then there are no main_votes for 1 in the next round */
forall (i in PROC) forall (r in ROUNDS) for(v = 0; v ≤ 1; v = v + 1) {
  inv3[i][r][v] : assert G ( (decide[r][i] ∧ decide_value[r][i]=v) ⇒
    (main_votes[r+1][v+1 mod 2]=0 ∧ ¬corrupted_main[r+1][v+1 mod 2] ) );
  forall (j in PROC) {
    subcase inv3[i][r][v][j] of inv3[i][r][v] for j=history_main_votes[r+1][v+1 mod 2];
    using (inv3[i][r][v][j]),
      inv2[i][r][v],

```

```

lemma9[r+1][0][0][M],
lemma9[r+1][0][1][M],
lemma9[r+1][1][0][M],
lemma9[r+1][1][1][M],
lemma25[r+1][v+1 mod 2],
/* abstractions */
VOTES⇒{0,M},
ROUNDS⇒{r,r+1},
/* free variables */
coin//free,
corrupted_main//free,
corrupted_main[r+1][v+1 mod 2],
corrupted_pre//free,
decide//free,
decide[r][i],
main//free,
main_votes//free,
main_votes[r+1][v+1 mod 2],
pc//free,
pc[j],
pre//free,
pre[r+1],
pre_proc//free,
pre_proc_votes//free,
pre_votes//free,
start//free
prove inv3[i][r][v][j];
}
}
/* if a party decides on 0, then there are no main_votes for abstain in the next round */
forall (i in PROC) forall (r in ROUNDS) for(v = 0; v ≤ 1; v = v + 1) {
  inv4[i][r][v] : assert G ( (decide[r][i] ∧ decide_value[r][i]=v) ⇒ ( main_votes[r+1][2]=0 ∧ ¬corrupted_main[r+1][2] ) );
  forall (j in PROC) {
    subcase inv4[i][r][v][j] of inv4[i][r][v] for j=history_main_votes[r+1][2];
    using (inv4[i][r][v][j]),
      inv2[i][r][v],
      lemma9[r+1][0][0][M],
      lemma9[r+1][1][0][M],
      lemma9[r+1][0][1][M],
      lemma9[r+1][1][1][M],
      lemma11[r+1],
      lemma16[r+1],
      lemma17[r+1],
      lemma25[r+1][2],
      corrupted2[r+1],
      corrupted6[r+1],
      /* abstractions */
      VOTES⇒{0,M},
      ROUNDS⇒{r,r+1},
      /* free variables */
      coin//free,
      corrupted_main//free,
      corrupted_main[r+1][2],
      corrupted_pre//free,
      decide//free,
      main//free,
      main_votes//free,
      main_votes[r+1][2],
      pc//free,
      pc[j],

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```

        pre//free,
        pre_proc//free,
        pre_proc_votes//free,
        pre_votes//free,
        round//free,
        round[j]
        prove inv4[i][r][v][j];
    }
}
/*-----AGREEMENT PROOF-----*/

```

/ if two parties decide in the same round they decide on the same value */*

```

forall (i in PROC) forall (r in ROUNDS) forall (j in PROC) {
    agree1[i][r][j] : assert G ( (decide[r][i]  $\wedge$  decide[r][j])  $\Rightarrow$  decide_value[r][j]=decide_value[r][i] );
    using inv1[i][r],
        lemma25[r],
        corrupted1[r],
        /* abstractions */
        VOTES $\Rightarrow$ {0,M},
        /* free variables */
        coin//free,
        corrupted_pre//free,
        corrupted_main//free,
        main//free,
        main_votes//free,
        pre//free,
        pre_proc//free,
        pre_proc_votes//free,
        pre_votes//free,
        start//free
        prove agree1[i][r][j];
}

```

/ if a party decides in round r and other party decides in round r+1, then they must decide on the same value */*

```

forall (i in PROC) forall (r in ROUNDS) forall (j in PROC) {
    agree2[i][r][j] : assert G ( (decide[r][i]  $\wedge$  decide[r+1][j])  $\Rightarrow$  decide_value[r+1][j]=decide_value[r][i] );
    using inv3[i][r],
        lemma25[r+1],
        corrupted1[r+1],
        /* abstractions */
        VOTES $\Rightarrow$ {0,M},
        ROUNDS $\Rightarrow$ {r,r+1},
        /* free variables */
        coin//free,
        corrupted_pre//free,
        corrupted_main//free,
        decide//free,
        decide[r][i],
        decide[r+1][j],
        main//free,
        main_votes//free,
        pre//free,
        pre_proc//free,
        pre_proc_votes//free,
        pre_votes//free,
        start//free
        prove agree2[i][r][j];
}

```

/ if party i decides in round r and is the first party to decide, */*
/ then any other party (j) must decide in round r or round r+1 */*

```

/* to prove this property we require some assumptions which we cannot prove because we are using ordset */

/* the premise is that no party has decided before round r */
/* noting that a party continues for one more round after it decides under fairness we have */
/* all parties reaches round r+1 */
/* and any party that does not decide in round r reaches round r+2 */
/* (note that therefore these assumptions includes fairness requirements) */
forall (i in PROC) forall (r in ROUNDS) {
  agreement_assumption1[i][r] : assert F ( round[i]=r+1 );
  agreement_assumption2[i][r] : assert G ( ¬decide[r][i] ) ⇒ F ( round[i]=r+2 );
  assume agreement_assumption1[i][r], agreement_assumption2[i][r];
}
/* the proof is then given by */
forall (i in PROC) forall (r in ROUNDS) forall (j in PROC) {
  agree3[i][r][j] : assert G ( (decide[r][i]) ⇒ F (decide[r][j] ∨ decide[r+1][j]) );
  using inv1[i][r],
    inv3[i][r],
    inv4[i][r],
    lemma5[r+1][M],
    lemma19[r][M],
    lemma19[r+1][M],
    lemma23[r][0][M],
    lemma23[r+1][0][M],
    lemma23[r][1][M],
    lemma23[r+1][1][M],
    lemma25[r],
    lemma25[r+1],
    lemma31[r+1],
    lemma32[r+1],
    lemma37[r],
    lemma37[r+1],
    corrupted1[r+1],
    corrupted5[r+1],
    coin2[r],
    coin2[r+1],
    /* assumed properties */
    agreement_assumption1[j][r],
    agreement_assumption2[j][r],
    /* abstractions */
    VOTES⇒{0,M},
    ROUNDS⇒{r-1,r,r+1},
    /* free variables */
    coin//free,
    corrupted_pre//free,
    corrupted_main//free,
    decide//free,
    decide[r][j],
    decide[r][i],
    decide[r+1][j],
    main//free,
    main_votes//free,
    pre//free,
    pre_proc//free,
    pre_proc_votes//free,
    pre_votes//free,
    round//free,
    round[j]
  prove agree3[i][r][j];
}
}
}

```